Containers in a File



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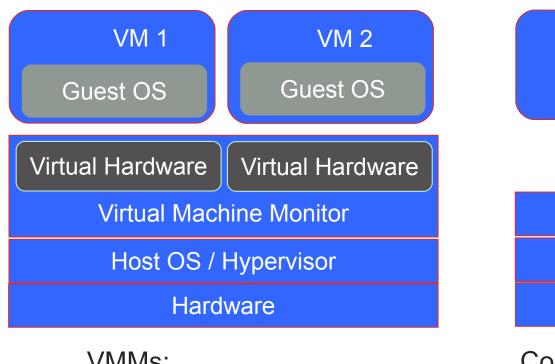
Agenda

Containers overview
Current way of keeping container files and its problems
Basics of container-in-a-file approach
Limitations of existing facilities
Solution
Benefits
Plans for future





Virtualization: VM-style vs. containers



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VMMs: ·KVM / Qemu ·VMware ·Parallels Hypervisor



OS Virtualization Layer

Host OS

Hardware

CT 1



CT 2

Containers Virtualization

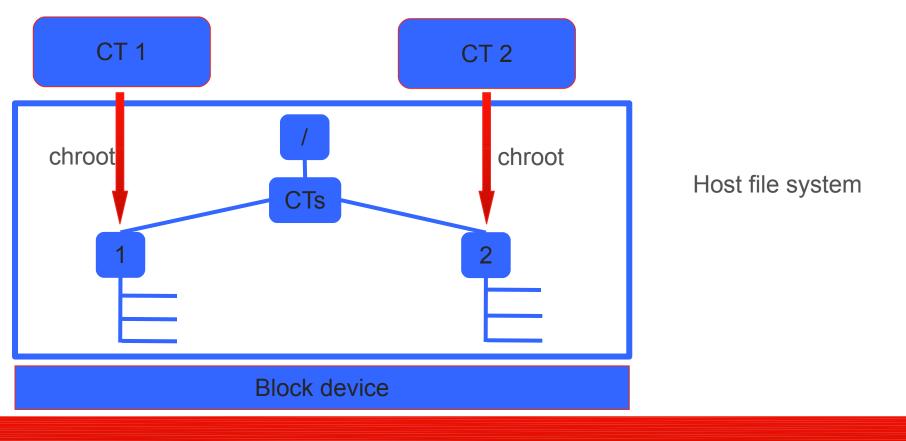
Each container has its own:

Files
Process tree
Network
Devices
IPC objects



Containers: file system view

Natural approach: each container chroot-s to a per-container sub-tree of some system-wide host file system.

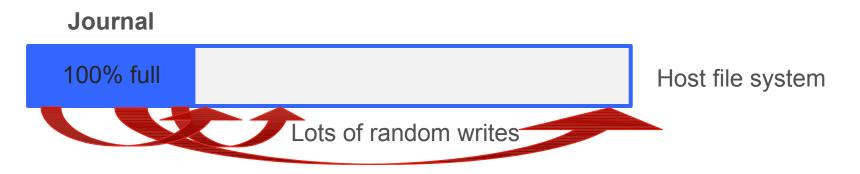




Problems

File system journal is a bottleneck:

1. CT1 performs a lot of operations leading to metadata updates (e.g. truncates). Consequently, the journal is getting full.



2. CT2 issuing a single I/O operation blocks until journal checkpoint is completed. Up to <u>15</u> seconds in some tests!



Problems (cont.)

Lots of small-size files I/O on any CT operation (backup, clone) Sub-tree disk quota support is absent in mainstream kernel

Hard to manage:
No per-container snapshots, live backup is problematic
Live migration – rsync unreliable, changed inode numbers

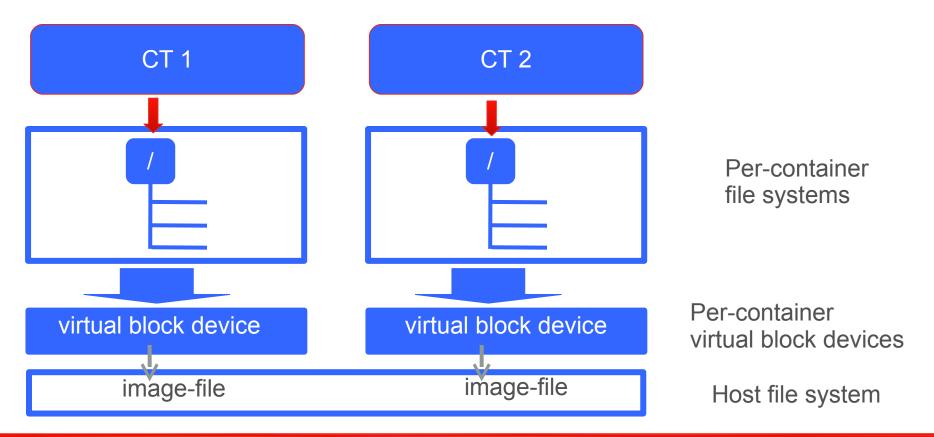
•File system type and its properties are fixed for all containers •Need to limit number of inodes per container





Container in a file

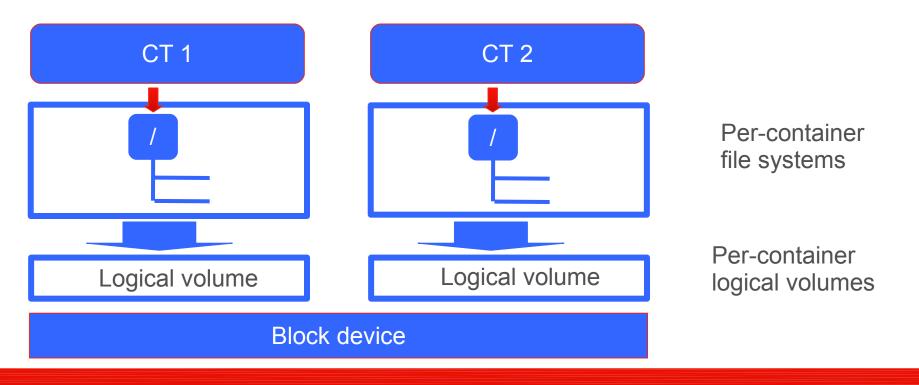
Basic idea: assign virtual block device to container, keep container' filesystem on top of virtual block device.





LVM limitations

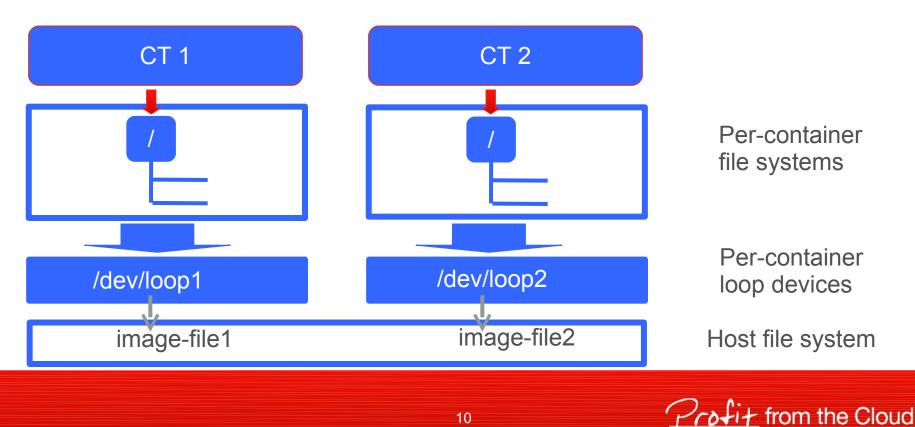
Not flexible enough – works only on top of block device
Hard to manage (e.g. how to migrate huge volume?)
No dynamic allocation (--virtualsize 16T --size 16M)
Management is not as convenient as for files



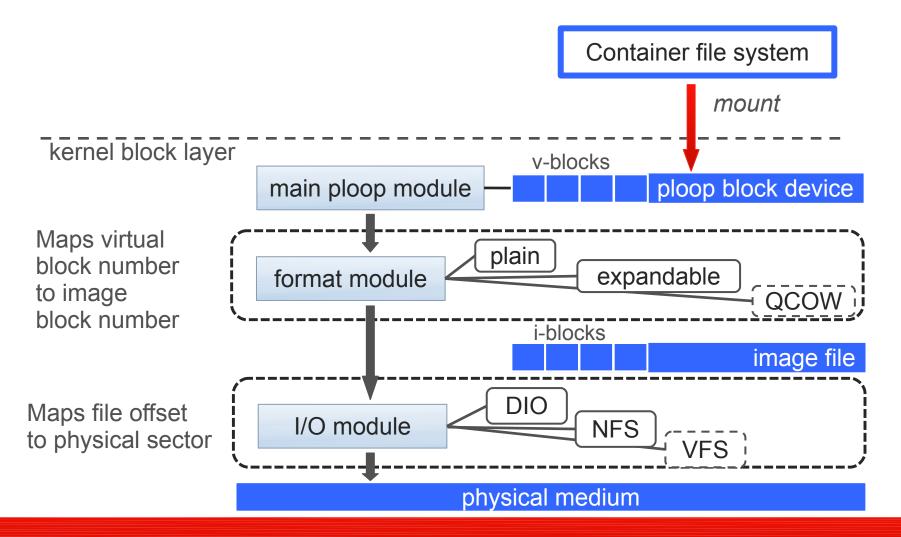


Ordinary loop device limitations

VFS operations leads to double page-caching
No dynamic allocation (only "plain" format is supported)
No helps to backup/migrate containers
No snapshot functionality



Design of new loop device (ploop)





Stacked image configuration

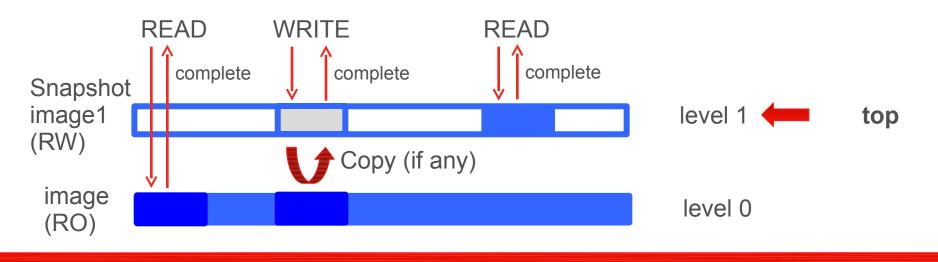
Ploop snapshots are represented by image files. They are stuffed in ploop device in stacked manner and obey the following rules:

Only top image is opened in RW mode (others – RO)

Every mapping bears info about 'level'

READ leads to access to an image of proper 'level'

WRITE may lead to transferring proper block to top image

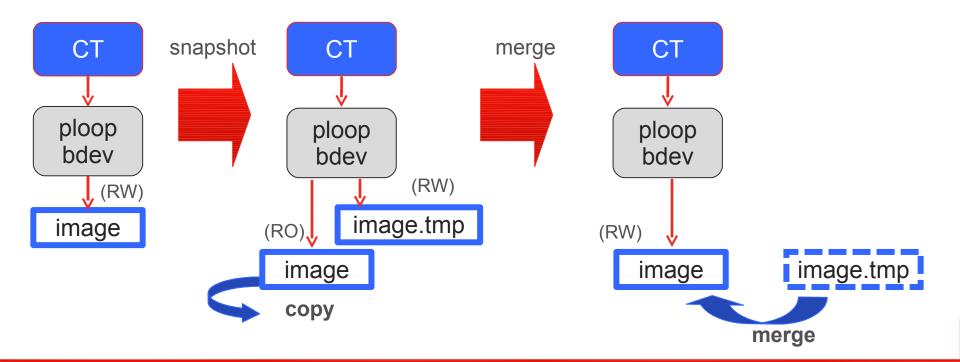




Backup via snapshot

Operations involved: •Create snapshot •Backup base image •Merge

Key features: •On-line •consistent

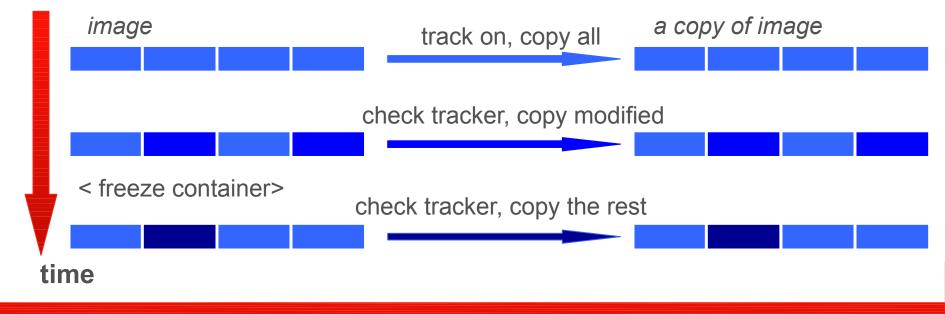




Migrate via write-tracker

Operations involved: •Turn write-tracker on •Iterative copy •Freeze container •Copy the rest

Key features: •On-line •I/O efficient



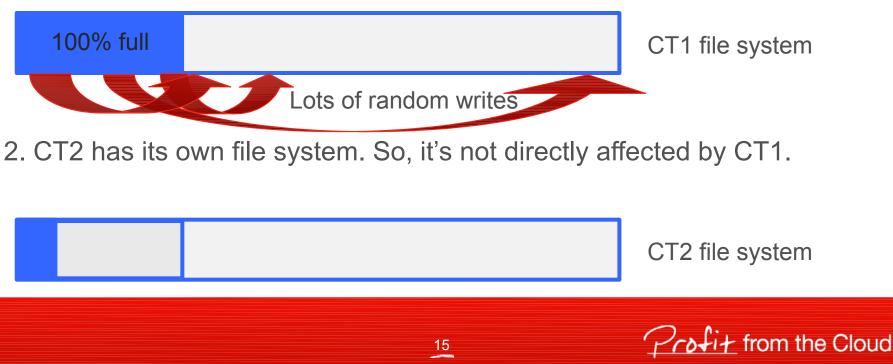


Problems solved

File system journal is not bottleneck anymore

1. CT1 performs a lot of operations leading to metadata updates (e.g. truncates). Consequently, the journal is getting full.

Journal



Problems solved (cont.)

Large-size image files I/O instead of lots of small-size files I/O on management operations

Disk quota can be implemented based on virtual device sizes. No need for sub-tree quotas

- Live backup is easy and consistent
- Live migration is reliable and efficient

Different containers may use file systems of different types and properties

No need to limit "number-of-inodes-per-container"



Additional benefits

Efficient container creation

Snapshot/merge feature

Can support QCOW2 and other image formats

Can support arbitrary storage types

Can help KVM to keep all disk I/O in-kernel







Implement I/O module on top of vfs ops

Add support of QCOW2 image format

Optimizations:
discard/TRIM events support
online images defragmentation
support sparse files

Integration into mainline kernel





Questions



